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Examination paper for
MA0301 Elementær diskret matematikk

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Problem 1

- a) How many non-negative integer solutions does the equation

$$x_1 + x_2 + x_3 + x_4 + x_5 = 21$$

have?

- b) How many of these solutions satisfy the additional requirements $x_2 \geq 2$ and $x_3 \geq 3$?
- c) How many of the solutions to (a) satisfy $x_4 \leq 4$ as an additional requirement?

Solution 1 a) There are $\binom{21+5-1}{21} = \binom{25}{21}$ non-negative integer solutions to this equation.

b) There are $\binom{16+5-1}{16} = \binom{20}{16}$ solutions to (a) that also satisfy $x_2 \geq 2$ and $x_3 \geq 3$.

c) There are $\binom{16+5-1}{16} = \binom{20}{16}$ solutions to (a) that satisfy $x_4 \geq 5$. Therefore, there are $\binom{25}{21} - \binom{20}{16}$ solutions to (a) that satisfy $x_4 \leq 4$.

Problem 2

- a) Determine whether or not the following statement is a tautology, by either proving it or giving a counterexample:

$$\left((p \rightarrow q) \wedge (r \rightarrow \neg q) \right) \rightarrow (p \wedge r)$$

- b) Negate the statement given in (a). (In your final answer, the \neg symbol may only appear directly in front of p , q and/or r .)
- c) Establish the validity of the following argument using the rules of inference

$$\begin{array}{l} p \rightarrow \neg q \\ \neg r \vee q \\ r \\ \hline \therefore \neg p \end{array}$$

Solution 2 a) Using the truth values $p = 0$, $q = 0$, $r = 0$, we find that $(p \rightarrow q)$ is true and $(r \rightarrow \neg q)$ is true. Therefore $((p \rightarrow q) \wedge (r \rightarrow \neg q))$ is true.

On the other hand, the statement $p \wedge r$ is false, and hence $((p \rightarrow q) \wedge (r \rightarrow \neg q)) \rightarrow (p \wedge r)$ is false. Therefore, the statement is not a tautology.

b) We negate the statement and work the negation through the brackets:

$$\begin{aligned} \neg\left(\left((p \rightarrow q) \wedge (r \rightarrow \neg q)\right) \rightarrow (p \wedge r)\right) &\iff \left((p \rightarrow q) \wedge (r \rightarrow \neg q)\right) \wedge \neg(p \wedge r) \\ &\iff \left((p \rightarrow q) \wedge (r \rightarrow \neg q)\right) \wedge (\neg p \vee \neg r) \end{aligned}$$

c) Using the rules of inference, we find that:

$$\begin{array}{ll} (1) & p \rightarrow \neg q & \text{premise} \\ (2) & \neg r \vee q & \text{premise} \\ (3) & r & \text{premise} \\ (4) & q & (2) + (3) + \text{rule of disjunctive syllogism} \\ (5) & \therefore \neg p & (1) + (4) + \text{modus tollens} \end{array}$$

Problem 3 Let $r \in \mathbb{R}$ with $r \neq 1$. Use induction to prove that

$$\sum_{i=0}^n r^i = \frac{1 - r^{n+1}}{1 - r}$$

for all $n \in \mathbb{Z}^+$.

Solution 3 We prove the formula by induction.

Basic step: For $n = 1$ we have

$$\sum_{i=0}^1 r^i = r^0 + r^1 = 1 + r$$

and

$$\frac{1 - r^{1+1}}{1 - r} = \frac{1 - r^2}{1 - r} = \frac{(1 + r)(1 - r)}{1 - r} = 1 + r.$$

Therefore, the formula holds for $n = 1$.

Now let $n \in \mathbb{Z}^+$ be arbitrary, and assume as an **Induction Hypothesis** that

$$\sum_{i=0}^n r^i = \frac{1 - r^{n+1}}{1 - r}.$$

Induction Step: Using the induction hypothesis, we now prove the formula for $n + 1$. We have

$$\begin{aligned} \sum_{i=0}^{n+1} r^i &= \left(\sum_{i=0}^n r^i \right) + r^{n+1} \\ &= \frac{1 - r^{n+1}}{1 - r} + r^{n+1} && \text{by the induction hypothesis} \\ &= \frac{1 - r^{n+1}}{1 - r} + \frac{(1 - r)r^{n+1}}{1 - r} \\ &= \frac{1 - r^{n+1}}{1 - r} + \frac{r^{n+1} - r^{n+2}}{1 - r} \\ &= \frac{1 - r^{n+1} + r^{n+1} - r^{n+2}}{1 - r} \\ &= \frac{1 - r^{n+2}}{1 - r} \\ &= \frac{1 - r^{(n+1)+1}}{1 - r} \end{aligned}$$

This shows that the formula also holds for $n + 1$, which completes our induction proof. The formula must therefore be valid for all $n \in \mathbb{Z}^+$.

Problem 4

- Give 6 different strings in the language $\{00\}\{101\}^* \cup \{011\}\{0\}^*\{01\}$
- Construct a finite state machine that recognises this language (with $\{0, 1\}$ as input and output alphabet).

Solution 4 a) There are many possible answers to this question. For example, take the strings 00, 00101, 00101101, 01101, 011001, 0110001.

- The Finite State Machine drawn in figure 1 recognises this language. There are also other valid solutions possible.

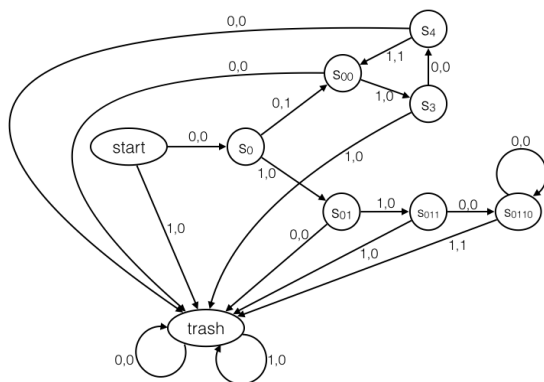


Figure 1: This figure belongs to Solution 4

Problem 5

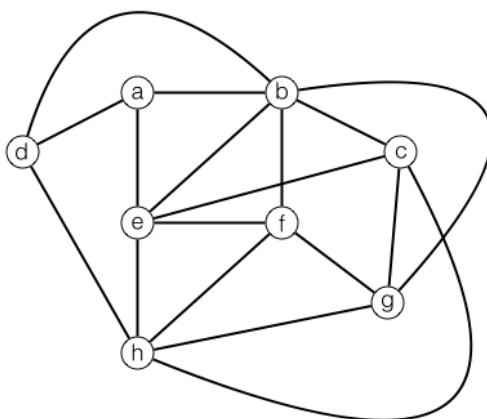


Figure 2: This figure belongs to Oppgave 5

- a) Find a rooted spanning tree with root a for this graph using a Breadth First Search Algorithm and vertex order a, b, c, d, e, f, g, h .
- b) Study the graph in Figure 1 and determine whether or not is it planar. Prove your answer.
- c) Does the graph in Figure 1 have an Euler Trail? Prove your answer.

Solution 5 a) We execute the Breadth First Search Algorithm with root a and vertex order a, b, c, d, e, f, g, h .

We start the tree with only vertex a , and queue $Q = (a)$. Next, we add vertices b , d and e to our tree with edges $\{a, b\}$, $\{a, d\}$ and $\{a, e\}$. Our new queue will be $Q = (b, d, e)$. We will now add c , f and g to our tree with edges $\{b, c\}$, $\{b, f\}$ and $\{b, g\}$. Our new queue will be $Q = (d, e, c, f, g)$. Finally, we add vertex h to our tree with edge $\{d, h\}$. Our queue will become $Q = (e, c, f, g, h)$. Visiting the vertices left in the queue will not give us any new vertices to add to our tree, since all vertices have already been added. See figure 3 for the BFS spanning tree.

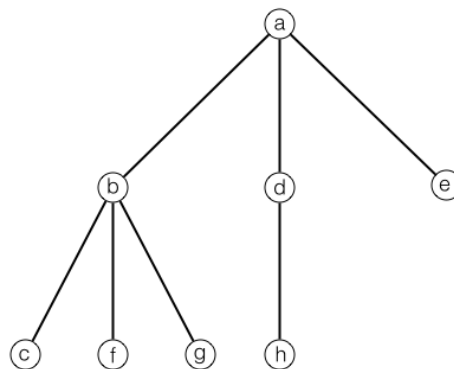


Figure 3: This figure belongs to Solution 5a

b) This graph is not planar. Consider the subgraph drawn in Figure 4, obtained by removing edges $\{h, f\}$, $\{b, f\}$, $\{a, b\}$, $\{a, d\}$, $\{a, e\}$ and vertex a . This graph can be obtained from K_5 by elementary subdivisions: draw K_5 and label its vertices b, c, e, g, h and apply an elementary subdivision to edges $\{e, g\}$ and $\{b, h\}$.

We see that our original graph contains a subgraph which is homeomorphic to K_5 . Therefore, it cannot be planar.

c) Notice that vertices d and a have degree 3, and vertices e and h have degree 5. Since the graph has more than two vertices of odd degree, it cannot have an Euler Trail.

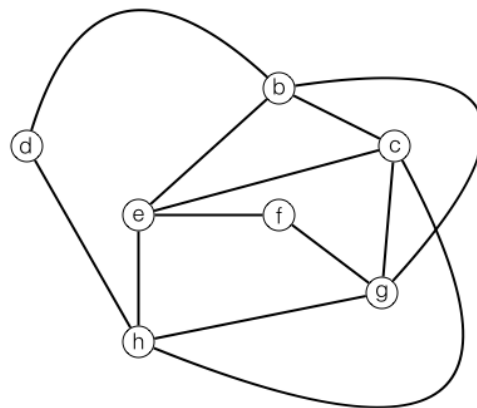


Figure 4: This figure belongs to Solution 5b

Problem 6 Let A be the set of all functions from \mathbb{Z}^+ to $\{1, 2, 3\}$.

- Give the three properties of an equivalence relation. You may give the names of these properties, or their definitions.
- Define a relation \mathcal{R}_1 on A by setting $f\mathcal{R}_1g$ if and only if $f(5) = g(5)$. Prove that \mathcal{R}_1 is an equivalence relation.
- Define a relation \mathcal{R}_2 on A by setting $f\mathcal{R}_2g$ if and only if there exists an $n \in \mathbb{Z}^+$ such that $f(n) = g(n)$. Prove that \mathcal{R}_2 is not an equivalence relation.

Solution 6 a) An equivalence relation \mathcal{R} on a set A has to be

- reflexive ($\forall a \in A : a\mathcal{R}a$)
- symmetric ($\forall a, b \in A : \text{if } a\mathcal{R}b, \text{ then } b\mathcal{R}a$), and
- transitive ($\forall a, b, c \in A : \text{if } a\mathcal{R}b \text{ and } b\mathcal{R}c, \text{ then } a\mathcal{R}c$).

b) We need to prove that \mathcal{R}_1 is reflexive, symmetric and transitive.

- Reflexivity: If $f \in A$, then naturally we have $f(5) = f(5)$. Hence $f\mathcal{R}_1f$. This proves reflexivity.
- Symmetry: If $f, g \in A$ and $f\mathcal{R}_1g$, then $f(5) = g(5)$. This implies that $g(5) = f(5)$, so $g\mathcal{R}_1f$. Therefore, the relation is symmetric.

- *Transitivity: If $f, g, h \in A$ with $f\mathcal{R}_1g$ and $g\mathcal{R}_1h$, then we know that $f(5) = g(5)$ and $g(5) = h(5)$. This implies that $f(5) = h(5)$, so $f\mathcal{R}_1h$. This proves transitivity.*

Since all three properties are satisfied, \mathcal{R}_1 is an equivalence relation.

c) *Define the function $f: \mathbb{Z}^+ \rightarrow \{1, 2, 3\}$ by $f(n) = 1$ for all $n \in \mathbb{Z}^+$.*

Define the function $g: \mathbb{Z}^+ \rightarrow \{1, 2, 3\}$ by $g(1) = 1$ and $g(n) = 2$ for all $n \geq 2$ with $n \in \mathbb{Z}^+$.

Define the function $h: \mathbb{Z}^+ \rightarrow \{1, 2, 3\}$ by $h(n) = 2$ for all $n \in \mathbb{Z}^+$.

Then f, g , and h are elements of A . Since $f(1) = 1 = g(1)$, we have $f\mathcal{R}_2g$. And as $g(2) = h(2)$, we have $g\mathcal{R}_2h$. Now note that for all $n \in \mathbb{Z}^+$ we have $f(n) = 1 \neq 2 = h(n)$, which implies that $f \not\mathcal{R}_2 h$. Therefore, the relation \mathcal{R}_2 is not transitive. This means that it cannot be an equivalence relation.