TTM4135 Information Security Exam June 6, 2005 --- solutions

1-25

Question		Choice		
	a	b	с	
1		$\checkmark$		
2	$\checkmark$			
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24		$\checkmark$		
25	✓			

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- 26. The recommendation X.509, part of the X.500 directory service recommendations, defines a format for public-key certificate that can be issued in a hierarchy of certification authorities. Figure 4.3 shows the formats.
- 27. 1) Verifiable signature 2) Within validity period.
- 28. Revocation because : 1) Change of authorization, for instance signer is not at the company anymore. 2) A private key is compromised or lost. 3) Certification system update, for instance algorithmic changes.
- 29. Invalidation by: 1) Time period expiration. or revocation list (CRL):2) Revocation of issuer's public key.3) Revocation of subject's public key.
- 30. See Section 5 of Security Lab Assignment, and Figure 4.4 in text book.
- 31. See textbook Figure 7.6.
- 32. Alice randomly selects two large prime number p and q, and computes n := p\*q. Then she randomly selects d, conditioned on that d is relative prime to phi(n) and 0 < d < phi(n). Alice must keep p, q and d secret.
- 33. Alice must provide Bob with her public key (e, n), where  $e := d^{-1} \mod phi(n)$ . Bob verifies by checking whether  $(s^{e} \mod n = m)$ .
- 34. If there are no restrictions on  $m^*$  then there are no restrictions on  $s^*$ . Bob starts with a random  $s^*$  and computes  $m^* := s^* e \mod n$ .
- 35. The hash function should satisfy: 1) Accept any length input. 2) Fixed length output. 3) 'Easy' to compute y := h(x). 4) Onewayness: 'Hard' to compute x given h(x). 5) Strong collision-free: 'Hard' to find a pair (x1, x2) such that h(x1) = h(x2).
- 36. Make h() public. In signing and verification replace *m* with h(m).
- 37. Alice computes E(K,m) and sends the result (c, t) to Bob. Bob decrypts the message *m* by computing  $c/t^k \mod p$ .

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